

Combining data-driven and symbolic reasoning for  
Invariant Synthesis in SMT  
(Work in Progress)

Haniel Barbosa

Andrew Reynolds

Cesare Tinelli



Clark Barrett

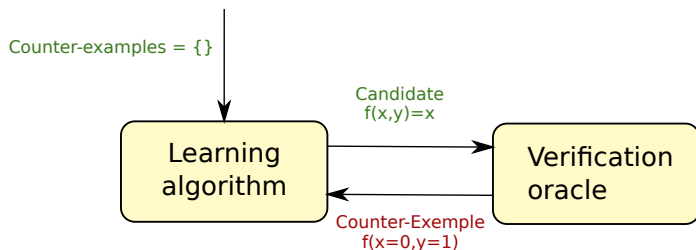


MVD 2018

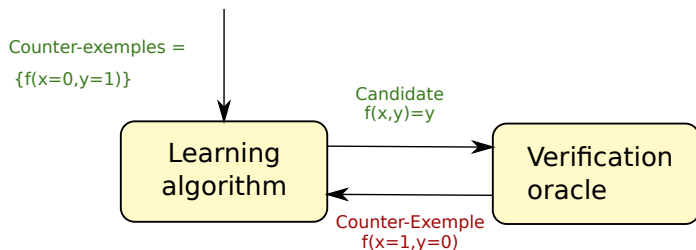
2018-09-29, Iowa City, IA, USA

## SyGuS Solving

- ▷ Most common technique for SyGuS solving
- ▷ Specification:  $x \leq f(x, y) \wedge y \leq f(x, y)$
- ▷ Expression search space:
  - ▶ Combinations of  $x, y, 0, 1, \leq, +$ , if-then-else



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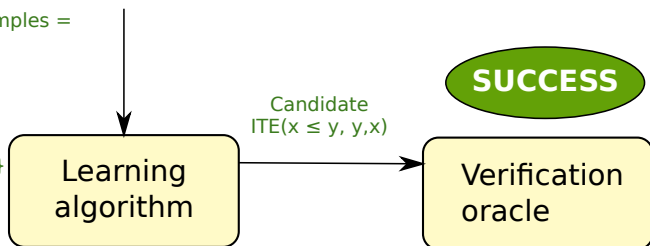
Counter-examples =

{ $f(x=0, y=1)$

$f(x=1, y=0)$

$f(x=0, y=0)$

$f(x=1, y=1)$ }



# Scalability issues

For this bit-vector grammar, enumerating

- ▷ Terms of size = 1 : .05 seconds
- ▷ Terms of size = 2 : .6 seconds
- ▷ Terms of size = 3 : 48 seconds
- ▷ Terms of size = 4 : 5.8 hours
- ▷ Terms of size = 5 : ??? (100+ days)

```
(synth-fun f ((s (BitVec 4))
              (t (BitVec 4))))
(BitVec 4) (
(Start (BitVec 4) (
  s t #x0
  (bvneg Start)
  (bvnot Start)
  (bvadd Start Start)
  (bvmul Start Start)
  (bvand Start Start)
  (bvlshr Start Start)
  (bvor Start Start)
  (bvshl Start Start))))))
```

- ▷ Generate partial solutions correct on subset of input
- ▷ Combine using conditionals

Step 1: Propose terms until all points **covered**

Step 2: **Generate predicates**

Partial Solutions

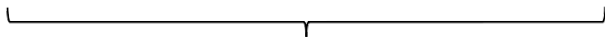
0  
1  
x  
y

Examples

(1, 1)  
(1, 2)  
(2, 1)  
⋮

Predicates

$0 \geq 1$   
 $1 \geq 1$   
 $x \geq 1$   
 $x \geq 2$   
 $x \geq y$



Step 3: **Combine!** *if ( $x \geq y$ ) then  $x$  else  $y$*

Only applicable for **plainly separable** specifications

A new framework for SyGuS solving



# CegisUnif : combining CEGIS with unification

- ▷ Not limited to plainly separable specifications
- ▷ *Data-driven*: refinement lemmas generate data points
- ▷ *Divide-and-conquer*: each point yields a new function to synthesize
  - ▶ Terms assigned to functions must satisfy refinement lemmas
  - ▶ SMT solving provides term candidates through constraint solving

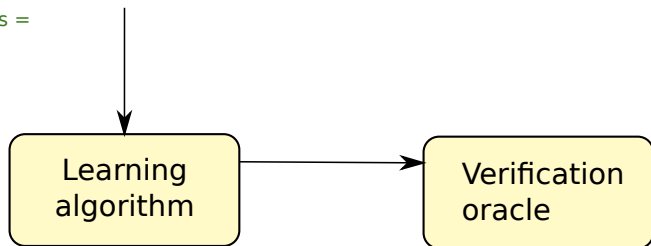
Counter-examples =

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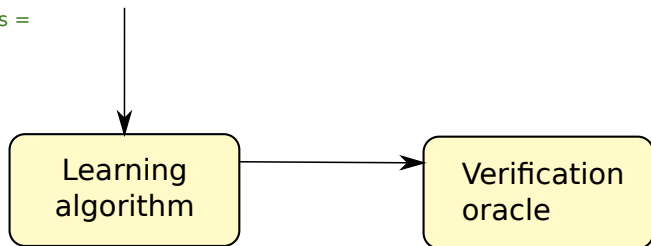
Counter-examples =

$f_1(x=0, y=1)$

$f_2(x=1, y=0)$

$f_3(x=0, y=0)$

$f_4(x=1, y=1)$



# Feature synthesis

- ▷ *Symbolic approach* : derive minimal number of features that separate conflicting points (i.e. those that cannot be assigned the same term)
  - ▶ Optimal fairness criteria?
    - Currently: consider terms of size up to  $\log_2(\#features)$
  
- ▷ *Heuristic approach* : accumulate “feature pool” and chose separating features based on information gain heuristic for decision tree learning
  - ▶ Select features that maximize information gain

Solving Invariant synthesis with CegisUnif

# Invariant Synthesis

```
Add(Int x, y) {  
  z := x; i := 0;  
  assume(y > 0);  
  while (i < y) {  
    z := z + 1;  
    i := i + 1;  
  }  
  return z;  
}
```

Post-condition:

$\forall x, y : z = x + y$

Result is the sum  
of the inputs

# Invariant Synthesis

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Invariant?

Post-condition:

$\forall x, y : z = x + y$

Result is the sum  
of the inputs

Verification:

$z = x \wedge i = 0 \wedge y > 0$	$\rightarrow$	$Inv(x, y, z, i)$
$Inv(x, y, z, i) \wedge i < y \wedge z' = z + 1 \wedge i' = i + 1$	$\rightarrow$	$Inv(x, y, z', i')$
$Inv(x, y, z, i) \wedge i \geq y$	$\rightarrow$	$z = x + y$

# Invariant Synthesis

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Add(Int x, y) {  
  z := x; i := 0;  
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    z := z + 1;  
    i := i + 1;  
  }  
  return z;  
}
```

$Inv(x, y, z, i)$   
 $z = x + i$   
 $z \leq x + y$

**Post-condition:**

$\forall x, y : z = x + y$

Result is the sum  
of the inputs

Verification:

$z = x \wedge i = 0 \wedge y > 0$	$\rightarrow$	$Inv(x, y, z, i)$
$Inv(x, y, z, i) \wedge i < y \wedge z' = z + 1 \wedge i' = i + 1$	$\rightarrow$	$Inv(x, y, z', i')$
$Inv(x, y, z, i) \wedge i \geq y$	$\rightarrow$	$z = x + y$

# Invariant Synthesis in SyGuS

- ▷ State-of-the-art: LoopInvGen [Padhi and Millstein 2017]: *data-driven* loop invariant inference with automatic feature synthesis
  - ▶ Precondition inference from sets of “good” and “bad” states
    - Feature synthesis for solving conflicts
  - ▶ PAC (*probably approximately correct*) algorithm for building candidate invariants
  
- ▷ “Bad” states are dependent on model of initial condition (no guaranteed convergence)
  
- ▷ No support for implication counterexamples



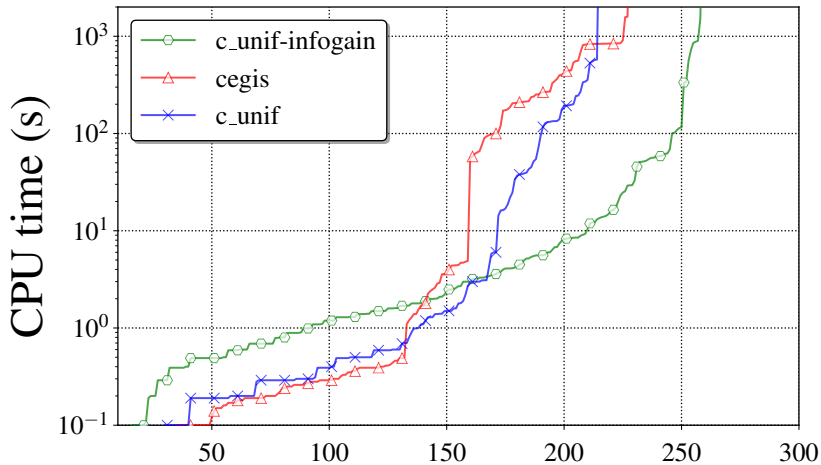
# Invariant Synthesis with CegisUnif

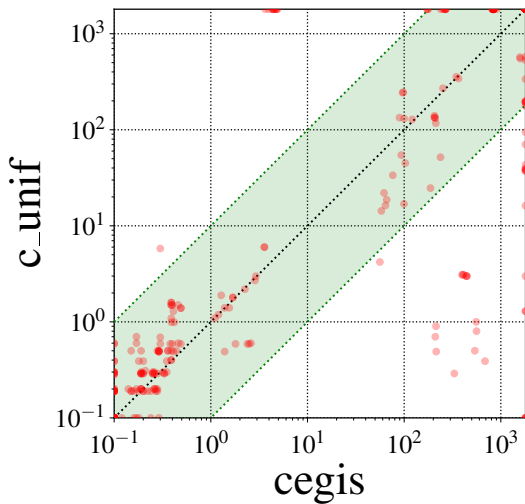
- ▷ Refinement lemmas allows derivation of three kinds on data points:
  - ▶ “good points” (invariant must always hold)
  - ▶ “bad points” (invariant can never hold)
  - ▶ “implication points” (if invariant holds in first point it must hold in second)
- ▷ No need for restriction to one initial state
- ▷ Native support for implication counterexamples
- ▷ Straightforward usage of classic information gain heuristic to build candidate solutions with decision tree learning
  - ▶ SMT solver “resolves” implication counterexample points as “good” and “bad”
  - ▶ Out-of-the-box Shannon entropy

Preliminary results

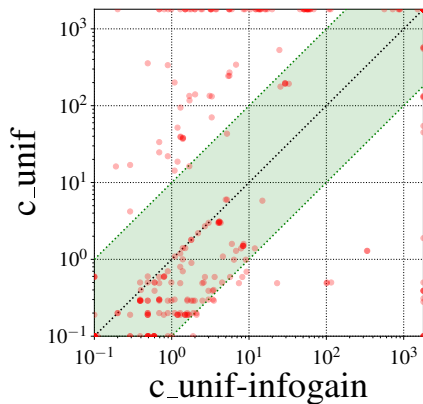
# Invariant generation for Lustre

- ▷ Test suite with 487 invariant synthesis benchmarks generated by the Kind 2 model checker from Lustre models
- ▷ We evaluate three configurations of CVC4
  - ▶ **cegis** : regular CEGIS
  - ▶ **c\_unif** : CegisUnif framework with symbolic solution building
  - ▶ **c\_unif-infogain** : CegisUnif framework with solution building determined by information gain heuristic
- ▷ 1800s timeout

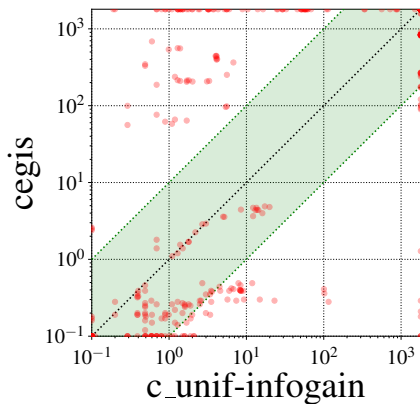




▷ + 38 / - 13



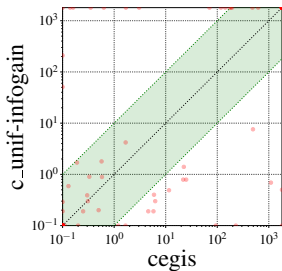
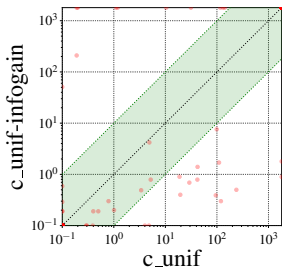
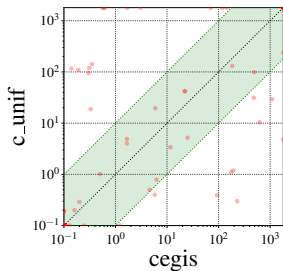
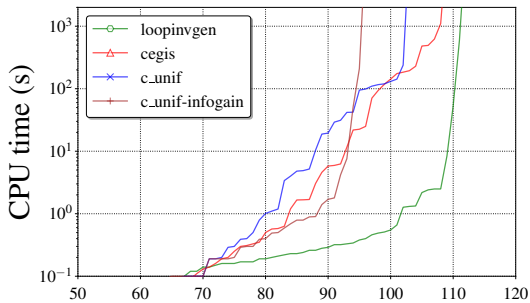
▷ + 63 / - 19



▷ + 73 / - 42

# Invariants category from SyGuS-Comp 2018

- ▷ Test suite with 127 invariant synthesis benchmarks from numerous applications
- ▷ We evaluate three configurations of CVC4
  - ▶ **cegis** : regular CEGIS
  - ▶ **c\_unif** : CegisUnif framework with symbolic solution building
  - ▶ **c\_unif-infogain** : CegisUnif framework with solution building determined by information gain heuristic
- ▷ We also compare against LoopInvGen, the current winner of the invariants category in SyGuS-Comp
- ▷ 1800s timeout





# Future work

- ▷ Adapt ICE [Garg et al. 2016] information gain heuristics to our setting; derive new heuristics
- ▷ Extend heuristics to function synthesis [Alur et al. 2017]
- ▷ Use data to determine “relevant arguments”
  - ▶  $f_1(0, 0, 0, 1, 2, 1, 0) \diamond f_2(1, 0, 0, 5, 2, 1, 3)$
  - ▶ Reducing noise: make points as similar as possible  
 $f'_1(1, 0, 0, 1, 2, 1, 0) \diamond f'_2(1, 0, 0, 5, 2, 1, 0)$
  - ▶ Only consider relevant arguments when synthesizing features
    - Can drastically reduce search space

# References

- Alur, Rajeev, Arjun Radhakrishna, and Abhishek Udupa (2017). “Scaling Enumerative Program Synthesis via Divide and Conquer”. In: Tools and Algorithms for Construction and Analysis of Systems (TACAS). Ed. by Axel Legay and Tiziana Margaria. Vol. 10205. Lecture Notes in Computer Science, pp. 319–336.
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- Padhi, Saswat and Todd D. Millstein (2017). “Data-Driven Loop Invariant Inference with Automatic Feature Synthesis”. In: CoRR abs/1707.02029. arXiv: 1707.02029.
- Solar-Lezama, Armando et al. (2006). “Combinatorial sketching for finite programs”. In: Architectural Support for Programming Languages and Operating Systems (ASPLOS). Ed. by John Paul Shen and Margaret Martonosi. ACM, pp. 404–415.